

# The Relational Model

Suan Lee

# Database Management System (DBMS)

- Used by all major commercial database systems
- Very simple model
- Query with high-level languages: simple yet expressive
- Efficient implementations

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### College


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### College

Name	State	ENR

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### Student

<b>ID</b>	<b>Name</b>	<b>GPA</b>	<b>Photo</b>
1	Adam	3.9	☺
2	Bill	3.4	<i>NULL</i>
3	Cathy	<i>NULL</i>	☺
⋮	⋮	⋮	⋮

### College

<b>Name</b>	<b>State</b>	<b>ENR</b>
Stanford	CA	15,000
Berkeley	CA	36,000
MIT	MA	10,000
⋮	⋮	⋮

- **Database** = set of named **relations** (or **tables**)
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- Each attribute has a **type** (or **domain**)

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- **Schema** = structural description of relations in database
- **Instance** = actual contents at given point in time

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- **NULL** – special value for “unknown” or “undefined”

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- **Key** – attribute whose value is unique in each tuple  
Or set of attributes whose combined values are unique

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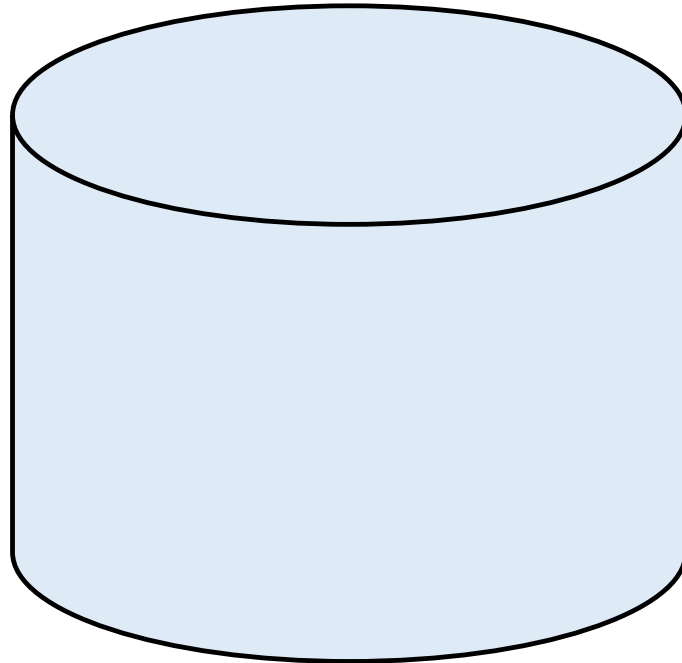
- Creating relations (tables) in SQL

```
CREATE TABLE Student (ID INTEGER, name STRING, GPA FLOAT, photo STRING)
```

```
CREATE TABLE college (name STRING, state CHAR(2), enrollment INTEGER)
```

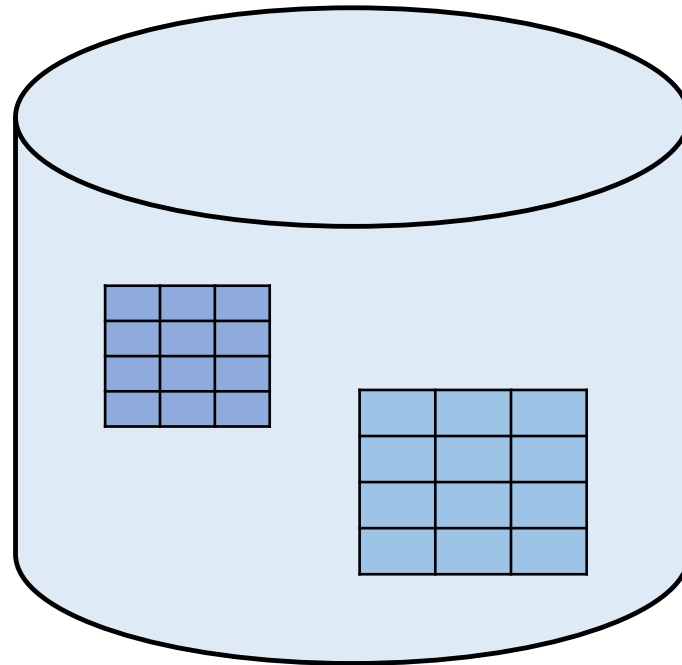
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# Steps in creating and using a (relational) database



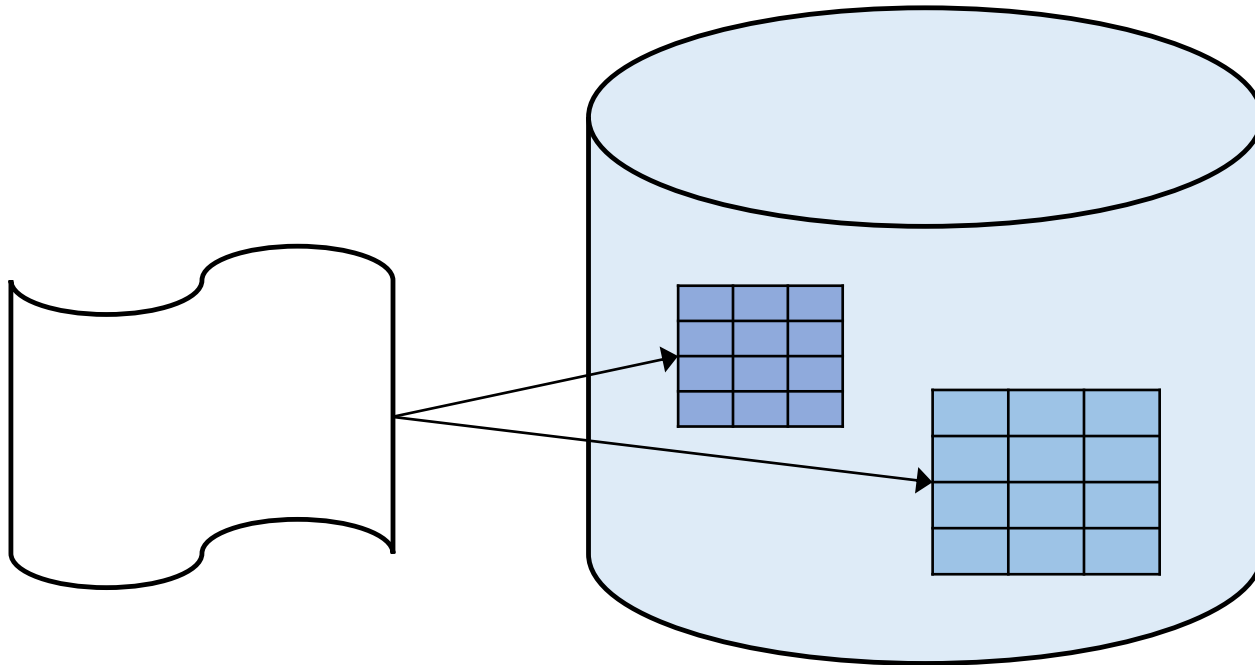
# Steps in creating and using a (relational) database

1. Design schema; create using DDL



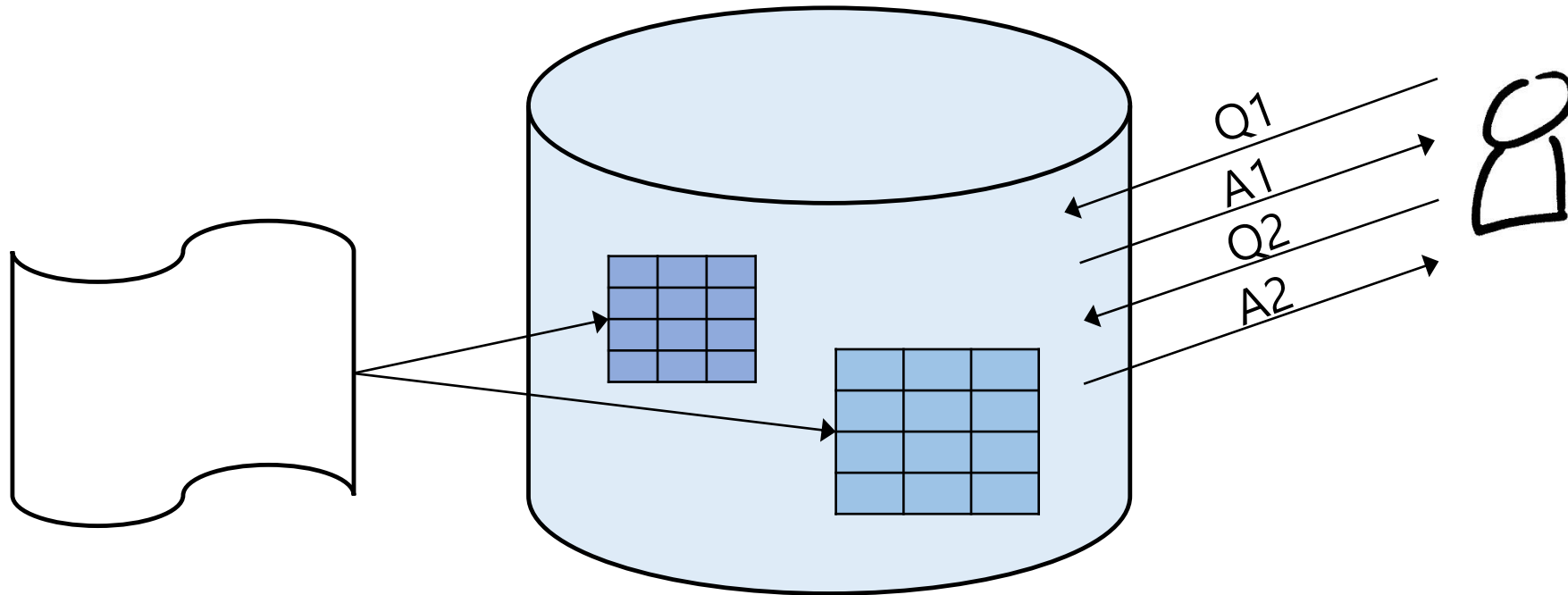
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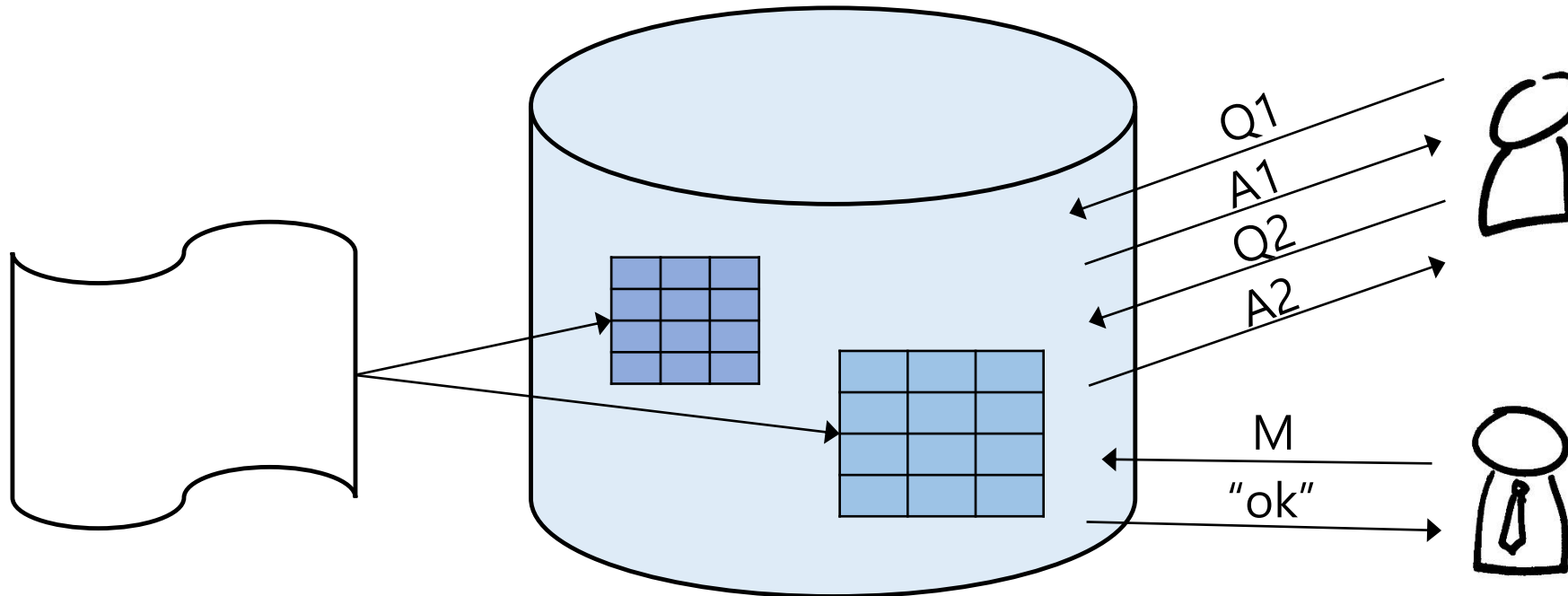
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3. Repeat: execute queries and modifications





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*All engineering departments in CA with < 500 applicants*

*College with highest average accept rate over last 5 years*

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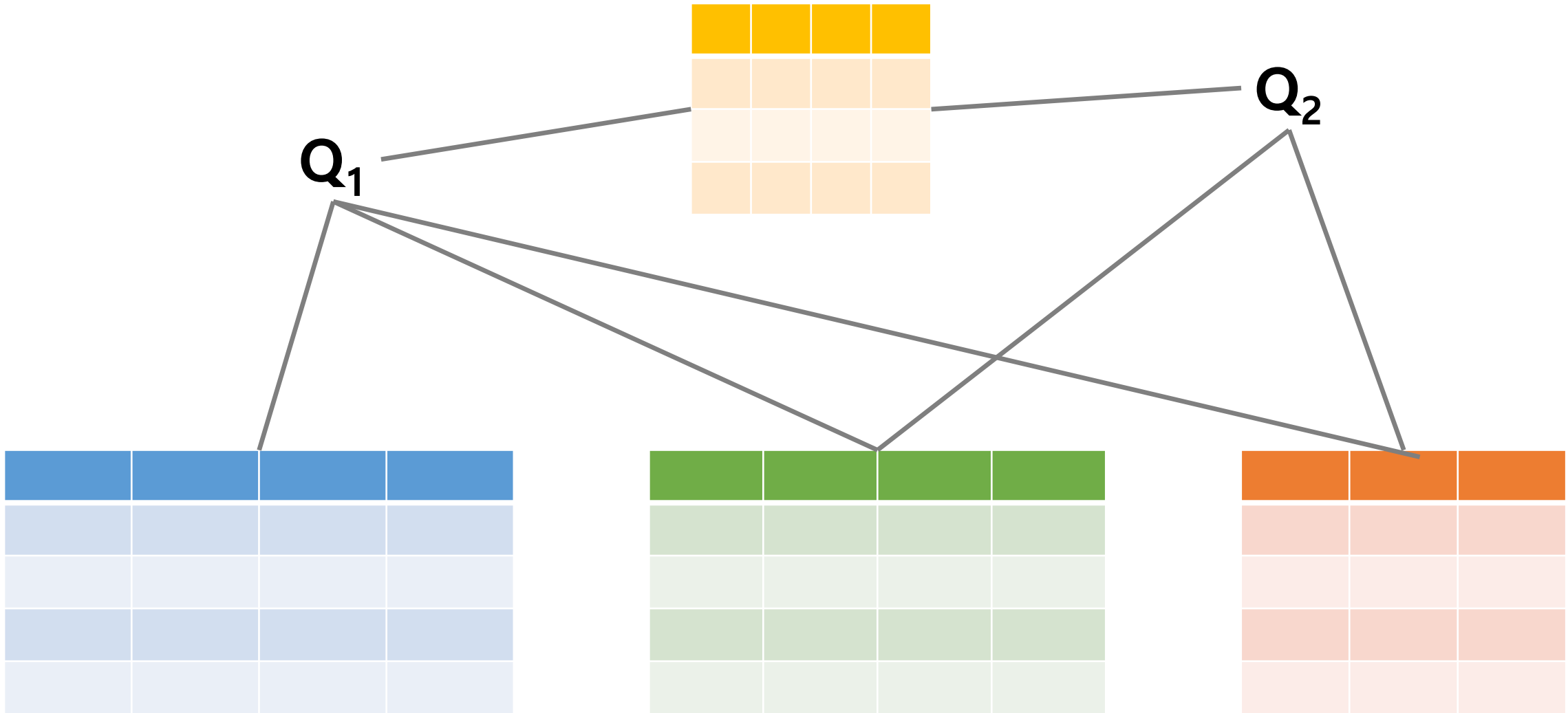
- Some easy to pose; some a bit harder
- Some easy for DBMS to execute efficiently; some harder
- “Query language” also used to modify data

# Queries return relations (“compositional”, “closed”)

Q<sub>1</sub>




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- Relational Algebra (formal)
  
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$$\Pi_{ID}(\sigma_{GPA=3.7 \wedge College.Name='Stanford'}(Student \bowtie Apply))$$

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$$\Pi_{ID}(\sigma_{GPA=3.7 \wedge \text{College.Name}='Stanford'}(\text{Student} \bowtie \text{Apply}))$$

- SQL (actual/implemented)

```
SELECT Student.ID  
FROM Student, Apply  
WHERE Student.ID=Apply.ID  
AND GPA>3.7 AND College.name='Stanford'
```

*IDs of students with GPA > 3.7 applying to Stanford*